



What's New for LMDB 1.0

Howard Chu CTO, Symas Corp. hyc@symas.com Chief Architect, OpenLDAP hyc@openIdap.org 2018-10-08

symas A Word About Symas

- Founded 1999
- Founders from Enterprise Software world
 - platinum Technology (Locus Computing)
 - IBM
- Howard joined OpenLDAP in 1999
 - One of the Core Team members
 - Appointed Chief Architect January 2007
- No debt, no VC investments: self-funded





Intro

- Howard Chu
 - Founder and CTO Symas Corp.
 - Developing Free/Open Source software since 1980s
 - GNU compiler toolchain, e.g. "gmake -j", etc.
 - Many other projects...
 - Worked for NASA/JPL, wrote software for Space Shuttle, etc.





Topics

(1) LMDB so far (0.9)(2) Coming in 1.0

symas (1) LDMB so far (0.9)

- Initially set out to replace BerkeleyDB in OpenLDAP alone, but now supported across multiple fields:
 - OS level rpm, git, SASL, Kerberos, Samba
 - Cryptocurrency blockchains
 - AI / deep learning / machine learning
 - High frequency trading
 - High performance computing

symas (1) LDMB so far (0.9)

- Released in 2011
 - Still the world's smallest, fastest and most reliable transactional embedded storage engine
- Now widely adopted
 - Dozens of wrappers for other languages (34 at last count)
 - Deployed as primary storage backend for countless open source and proprietary projects





(2) Coming in 1.0

- Support for 64bit DBs on 32bit builds (VL32)
- Page level encryption and/or checksums
- Incremental backup
- Long key support (up to 2GB keys)
- Semi-synchronous writes





(2) Coming in 1.0

- Headerless overflow pages
- Restructured freelist
- Raw partition support
- 2-phase commit





(2) Coming in 1.0

- On-disk format change
 - Most of these features require a change to the LMDB page formats
 - Some of these features have been working for years, but we didn't want to release them separately and have to deal with multiple DB migrations



(2) VL32



- Optional support for larger-than-32bit databases on 32bit builds
 - Instead of using a single mmap for the entire DB, maps in chunks of >= 16 pages at a time
 - Requires some user-level caching, slower than native 32bit support
 - Must track referenced pages per transaction and within the entire environment
- Allows DBs to be used interchangeably between 32bit and 64bit machines
 - If using semaphores instead of shared mutexes, DBs can be shared by 32bit and 64bit processes on the same machine





(2) Page encryption

- (Never thought we'd ever do this)
 - Requires app-level page caching
 - to keep a decrypted copy of a page being accessed
 - But we already have this, due to VL32 support
- Ciphers in user-supplied callbacks
 - there will be no built in encryption or checksum algorithms
- Needs space for per-page checksums

symas (2) Incremental Backup

- Adds transaction ID to page header
 - Allows simple backup of all pages newer than a specified txnID
 - Incremental backups can easily be overlaid on top of a snapshot from the specified txnID





(2) Long Keys

- Support for storing keys on overflow pages
 - Key length up to 2^31 -1 bytes
 - Values still limited to 2^32 -1 bytes





- Uses 2 new temporary meta pages for recording the semisync commits
- Use explicit mdb_env_sync() to persist temporary meta pages onto main meta pages
- In case of crash, un-synced commits will be lost, but main meta pages will be intact
- Will incur additional DB growth while un-synced





(2) Headerless OVF

- Currently the first page in a span of overflow pages always has a standard page header (16 bytes on 64bit build)
- So, storing a page-size value consumes two pages, because it can't all fit on the first page
- Relevant header info will be moved into the leaf node instead
 - page number, transaction ID, checksum





(2) Freelist

- Currently the freelist is stored as a single record per transaction
 - performance suffers when the free space is fragmented
- Instead, will use DUPSORT|DUPFIXED to let LMDB store each freelist entry in arbitrary pages
 - Total space requirement will be about the same
 - But will be immune to fragmentation issues





(2) Raw partitions

- Allow storing LMDB data directly on a raw block or character device
 - Avoids all filesystem metadata overhead
 - Bulk load on raw device is 2x faster than on ext4fs with journaling disabled
 - No other performance impact
 - Makes maxsize discussion totally moot
 - Allows exploiting special features of NVDIMMs



- Add an mdb_txn_prepare() API
- Enables support for distributed transactions
 - For use with slapd back-ldap / back-meta / etc.





Questions?